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A Theory of Commerce

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Abstract

Retail trade absorbs vast scarce resources because the physical process of trading is time consuming, buyers match with sellers without coordination, and consumers prefer to purchase a diverse basket of goods. Sellers post prices to attract customers, but buyers care also about the expected time it takes to make a purchase. Retail prices can be non-linear due to packaging and quantity discounts. However, prices cannot depend on buyers preferences because these are private information. To capture these features, we adopt directed search and assume that sellers ignore their clients preferences. If, realistically, sellers cannot charge a flat fee to all potential buyers, then in equilibrium the average lineup of buyers in front of a seller is inefficiently long. In contrast, the directed search equilibrium is efficient with full information. Our model is easily inserted in a Neoclassical growth framework. The retail trade sector can be calibrated using commercial margins and resources employed in that sector.

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1 Introduction

Commercial margins constitute a very large fraction of the cost of purchasing a typical good. On average, for every dollar spent before tax in a retail store in the United States, only 57 cents reach the producers of the merchandise. The remaining 43 cents cover the margins of the merchant wholesaler (15 cents) and the retailer (28 cents).¹ Models with Walrasian markets abstract from commercial activities. This abstraction would be of little importance if these activities were limited to transporting commodities from factories to consumers, for then we could think of commercial costs are incurred in the process of contacting and dealing with buyers in bilateral meetings. Moreover, the size of these costs is independent of the quantities that buyers end up acquiring and cannot be recovered if the buyer chooses to buy nothing. These facts can only be properly modeled with non-Walrasian markets. To develop an appropriate framework of analysis, we advance a directed search model of commerce. This model incorporates the multidimensional competition we observe in retail trade where buyers and sellers care not only about prices but also about the expected time it takes to perform a trade.² Moreover, the model incorporates the fact that sellers cannot directly observe the buyers' willingness to pay.

Even though the purchase of a single item appears to involve minimal costs, the retail trade sector absorbs vast scarce resources for three main reasons: First, the physical trading process is time consuming and requires proper facilities, commercial inventories, and other inputs. Second, buyers do not coordinate their decisions as to when and where to go shopping, so quite often some buyers are lining up for assistance in front of a seller, while another seller is idle. Third, consumers are willing to sacrifice some of their time to purchase a diverse basket of goods.

To capture these features in a simple model, we consider an environment where buying and selling are separate activities which require some of the scarce time available in the economy. As explained below, traders can direct their search to a segment of the market, but inside this segment

¹This break-up assumes that the good goes through one merchant wholesaler and one retailer. In our calculations, we use the gross commercial margins over sales ratios calculated by the Bureau of the Census, Service Sector Statistics Division for 2000 (0.278 for retail and 0.2054 for wholesale). This data was downloaded in July 1, 2002 from http://www.census.gov/svsd/www/artstbl.html and http://www.census.gov/svsd/www/atspur.txt.

²The concept of directed search was introduced by Peters (1991) and Montgomery (1991). While these authors assume a finite number of traders, we find it convenient to simplify strategic considerations by assuming a large number of traders as in Moen (1997) and Shimer (1996). The concept is then sometimes referred to as competitive search. See also Mortensen and Wright (2002).

buyers and sellers meet randomly in bilateral pairs according to a matching technology as in the search model of Mortensen (1982) and Pissarides (1990). Finally, due to a technological constraint, sellers specialize in selling a single good. Therefore, households face a trade-off between the diversity of goods they purchase and the time they spend shopping. In equilibrium, they choose to buy many goods because of their preference for diversity. Although our model provides a natural framework to address the endogenous creation of stores which sell many goods simultaneously, the holding of commercial inventories, and the inclusion of fixed capital in retail trade, we leave these extensions for future work.

Competition among retailers is keen. Sellers post and commit to prices to attract customers to visit their outlets. However, buyers care not only about prices. They also care about the expected time it takes for them to complete a purchase, or the probability of making a purchase in a given time interval. Since matching between buyers and sellers is uncoordinated, the probability of making a purchase in a given time interval depends on the average lineup of buyers in front of a seller (the ratio of buyers to sellers).

To model the combination of random matching and competition that we observe in retail trade, we assume that sellers post prices, and buyers direct their search towards the subset of sellers announcing a particular price offer. We refer to the set of sellers announcing a particular price offer and the buyers that direct their search towards them as a submarket. The assignment of buyers to sellers in a submarket is dictated by an exogenous matching technology. As noted by Acemoglu and Shimer (1999), this framework encompasses many reasonable possibilities. For example, it includes the possibility that symmetric buyers use identical mixed strategies to select the seller they visit as in the frictional assignment literature (Peters (1991), Montgomery (1991), Burdett, Shi, and Wright (2001), and Peters (2000)). Another possibility is that submarkets are located in separate places (malls or streets) and matching in each submarket takes place according to a matching technology.

As in Moen (1997) and Shimer (1996), we find that a directed search equilibrium with full information is efficient. However, we find that the qualifications to this efficiency result are more interesting than the result itself. In retail trade, each seller potentially serves a variety of buyers, each one with a particular willingness to pay. We refer to this willingness to pay as the buyer's type. With many types purchasing in the same submarket, pricing in a directed search equilibrium with full information must have one of two unrealistic properties. Either prices must be type dependent, or they must include a flat fee that all potential buyers must pay to cover retail costs.

The type dependent prices necessary to implement the direct search equilibrium with full infor-

mation requires that preferences be directly observable. This is clearly unrealistic and trivializes the complexities of retail trade pricing. In our baseline model, we assume that buyers' preferences are private information, so prices cannot depend on them. As we show, this is equivalent to assuming that in equilibrium bilateral trades between buyers and sellers must be incentive compatible, in the sense that buyers have no incentive to misrepresent their preferences.³

The second way to implement the directed search equilibrium with full information is by introducing flat fees. In retail trade, we observe some instances in which flat fees are charged to buyers prior to their purchasing decision. For example, some warehouse clubs charge membership fees. However, these fees are the exception rather than the rule. Moreover, even when they are present, they cover only a small fraction of retail costs. Explaining why flat fees are not common practice is an open question of upmost importance. We do not give a final answer to this question. Rather we construct a stylized model of retail trade where flat fees cannot arise as an equilibrium outcome. In our baseline model, buyers learn their valuations as soon as they meet a seller and observe the good the seller carries. This assumption, combined with the fact buyers' valuations can be arbitrarily low, rules out flat fees. Buyers with small valuations will simply refuse to pay the fee.⁴ However, our environment does not rule out other forms of non-linear pricing. In fact, the model predicts that per unit prices should decline with quantity. This may sound counterfactual. However, retailers use both packaging and explicit quantity discounts to offer these decreasing per unit prices. For instance, when we purchase a gallon of paint we pay much less than when we purchase four separate quarters of a gallon of the same paint.

With the informational environment of our baseline model, the directed search equilibrium is not efficient. In equilibrium the average lineup of buyers in front of a seller is inefficiently long. Nevertheless, a policy maker who faces the same informational constraints as the sellers and whose only policy instrument is to regulate market prices cannot improve upon the equilibrium allocation when the production technology is linear (affine production function).⁵ Thus, if the production

³Admittedly, some transactions in retail trade involve bargaining, and the bargaining process may partially reveal the willingness to pay of a buyer. Yet, the amount of information revealed in the bargaining process is also limited by incentive compatibility constraints. Indeed, in the environment we study in this paper there is no exogenous bargaining rule that attains the efficient outcome when buyers' preferences are unobservable. Moreover, bargaining in retail trade is rare and is inconsistent with ex-ante price posting.

⁴Using standard terminology in the mechanism design literature, pricing mechanisms which include flat fees violate the individual rationality constraints of the buyers. On the other hand, the type dependent prices necessary to implement an efficient allocation violate the buyers' incentive compatibility constraints.

⁵In general equilibrium, efficiency in the commercial sector is not enough. We also need production efficiency.

technology is linear, the directed search equilibrium is second best efficient.⁶ Assuming a linear technology is certainly restrictive. However, as long as the production function is continuously differentiable, we can approximate it with an affine production function. Therefore, as long as the policy maker faces the same informational constraints as the sellers, the welfare gains of regulating market prices are second order of magnitude. We use numerical simulations to check that these potential welfare gains are minimal.

Our model can be easily incorporated in a Neoclassical growth framework with two sectors. One of these sectors produces goods combining capital and labor as is typical in the Neoclassical framework, while the other exchanges goods in retail markets. The combined model can be estimated using standard empirical data. In particular, the parameters of the retail sector can be identified and estimated using commercial margins and the average time allocation of households.⁷ The tractability of the model makes it suitable for further extensions. For example, one can use this framework to introduce money along the lines of the search theoretical approach of Kiyotaki and Wright (1989 and 1993). Faig (2001) constitutes a first attempt in this direction in a simpler version of the present model. In this simpler version, sellers are restricted to make offers to the buyers they are paired with which consist of a single quantity-payment pair. Our main improvement with respect to Faig (2001) is the relaxation of this constraint by allowing sellers to make offers that consist of a price schedule that maps the quantity chosen by a buyer into the payment to be made to the seller. This price schedule serves as a mechanism to reveal the buyers' private information about their preferences. Also, in Faig (2001) search is undirected and sellers make take-it-or-leave-it offers to buyers.

Two recent papers, Soller-Curtis and Wright (2000) and Camera and Delacroix (2001) also study search-theoretic models where the buyers' willingness to pay for a good is private information. In both of these papers, goods are indivisible and search is undirected. Soller-Curtis and Wright assume that sellers make take-it-or-leave-it offers to buyers and they focus on the coexistence of two prices for the same good in equilibrium. Camera and Delacroix focus on the endogenous determination of the trade mechanism - sellers can choose if they want to commit to a posted With a linear technology, the marginal return to labor in the production sector is fixed. Thus, the global allocation of labor is optimal as long as buyers and sellers are given the right incentives in the commercial sector.

⁶With full information, the restriction to linear technologies is not necessary.

⁷Shi (1999 and 2001a) also incorporates a search model in a Neoclassical growth framework. The model in Shi does not differentiate between producers and sellers as we do. Also, his model assumes full information and random (undirected) search.

price or to bargain once they meet a buyer. When buyers have identical preferences and search is undirected, pre-committing to a price is preferable from the sellers point of view because it allows them to extract the whole trading surplus. However, when the preferences of the buyers are heterogeneous, the bargaining process allows the seller to infer information about the buyers' preferences and hence to price discriminate. In our paper, goods are divisible, so sellers can commit to a non-linear price schedule which allows a restricted form of price discrimination even without bargaining. Also, we incorporate the equilibrium concept of directed search which endogeneizes the market power of sellers in a reasonable fashion.

Peters and Severinov (1997) also study a model of directed search where the preferences of the buyers are private information. However, in their model the price setting mechanism is an auction among the buyers that meet a seller. Moreover, because a single unit is sold in each auction, their paper does not deal with the price incentives to endogenously determine the size of each purchase.

The paper is organized as follows. Section 2 presents an overview of the model and our equilibrium concept. Section 3 characterizes the optimal behavior of a representative household for given prices. Section 4 studies how prices are endogenously determined under directed search by analyzing the interaction between buyers and sellers in the market place. Section 5 combines the analyses of the previous two sections into a general equilibrium model where both the behavior of households and prices in retail trade are endogenous. Section 6 studies the welfare properties of a directed search equilibrium. Section 7 incorporates a generalized version of our model in a dynamic Neoclassical growth framework and discusses how to identify the parameters using standard data. Section 8 briefly discusses some of the issues one must confront when extending the present model and concludes. The proofs are gathered in the Appendix.

2 Overview of the Model

The economy consists of a continuum of households with measure one who produce and consume differentiated goods. Households do not consume the goods they produce so they need to trade. Trading activities involve some degree of idiosyncratic uncertainty to be specified below. To avoid the ex-post heterogeneity induced by idiosyncratic uncertainty, which severely limits the tractability of the model, we follow Shi (1997) in assuming that each household is composed of a large number of individuals. These individuals independently perform the production and exchange activities in the household. Thanks to the large household assumption, each household faces no uncertainty, even

though the members of the household who perform trading activities are subject to idiosyncratic risks. All members of the household equally share the consumption of goods which is their only source of utility, so there is no conflict of interests among them.⁸

To construct a simple environment where households buy the goods they consume in a retail trade sector, we make the following assumptions. Each household produces a single divisible good and likes to consume the goods produced by other households. Because of physical constraints, the members of the household who perform manufacturing activities (producers) cannot simultaneously sell the good they produce. Likewise, sellers must remain in their retail outlets to sell their products. Therefore, for trade to take place the buyers of a household must go around visiting sellers of other households. Since sellers never meet each other, direct barter is ruled out.

All payments are denominated in an abstract numeraire. In the version of the model analyzed in this paper, buyers do not need to carry money with them. Instead, all traders have access to a central clearing-house that records the credits (payments received by sellers) and debits (payments made by buyers) of all households and ensures that their budget constraint is satisfied.

Trading activities are subject to two kinds of uncertainty. First, buyers sometimes find goods that, because of idiosyncratic factors, fit well the needs of their households while in other occasions they do not. Second, because of matching frictions, trading meetings between buyers and sellers are partially random. Thus, a trader may or may not be able to perform a transaction during a given period.

We model the first type of uncertainty by assuming a preference shock ε that scales the utility that a good brings to the household. Preference shocks are realized once a buyer of the household meets a seller. The realized value of ε is the buyer's own private information and cannot be observed by the seller. Therefore, in a trading meeting the seller ignores the willingness to pay of the buyer. We believe that this ignorance is key for understanding retail pricing in the same way that unobservable characteristics of taxpayers are key for understanding income tax schedules.

In order to model both the matching frictions and the price competition that characterize retail markets, we assume the following form of directed search.⁹ Prior to the trading period, each seller j simultaneously posts and commits to a price schedule $\mathcal{Z}_j(q)$. This schedule specifies the payment required in a transaction as a function of the quantity exchanged.¹⁰ In the next stage,

⁸The absence of a conflict of interests between the household and its members is an important difference between this paper and Shi (1997).

⁹Our description of this concept parallels the formulation in Acemoglu and Shimer (1999).

 $^{^{10}}$ Equivalently, because each seller j may meet buyers with different private valuations ex post, the seller announces

buyers observe the menu of price schedules posted by all sellers and simultaneously choose where to shop. That is, each buyer *i* decides to trade at a particular price schedule $\mathcal{Z}_i(q)$ in the set of posted schedules $\{\mathcal{Z}_j(q) : \text{for all } j\}$.¹¹ If buyer *i* meets a seller posting $\mathcal{Z}_i(q)$ during the trading period, the buyer privately observes his valuation ε for the seller's good, chooses the quantity to be transacted and pays for it according to the posted price schedule. Buyers can always choose not to buy anything in which case they pay nothing. If no meeting takes place, then there is no trade.¹² We refer to the set of sellers posting $\mathcal{Z}_j(q)$ and the set of buyers that direct their search to this price schedule as submarket *j*.

Depending on buyers's search decisions, there may be longer lineups in some submarkets than in some others. To capture this, we let $\theta_j \in [0, \infty]$ be the ratio of buyers to sellers in the submarket with a posted price schedule $\mathcal{Z}_j(q)$:

$$\theta_j = \frac{B_j}{S_j},\tag{1}$$

where S_j is the measure of sellers posting $\mathcal{Z}_j(q)$ and B_j is the measure of buyers who decide to trade at this price schedule. We refer to θ_j as the congestion in submarket j, or the average queue of buyers in front of a seller in this submarket. For the time being and to facilitate the exposition, we assume that buyers and sellers may perform at most one trade during the trading period (Subsection 5.1 relaxes this assumption). In a submarket with congestion θ_j , the probability that a seller meets a buyer is $m^s(\theta_j)$, where $m^s : [0, \infty] \to [0, 1]$ is continuously differentiable, decreasing, and concave. Symmetrically, a buyer meets a seller with probability $m^b(\theta_j)$, where $m^b : [0, \infty] \to [0, 1]$ is continuously differentiable, increasing, and convex. If many buyers seek a few sellers (θ_j is high), then it is easy for a seller to find a buyer and hard for a buyer to find a seller. By having m^s and m^b depend only on θ_j , we implicitly assume constant returns in matching,¹³ so

$$\mathcal{M}(B_j, S_j) = B_j m^b(\theta_j) = S_j m^s(\theta_j), \tag{2}$$

a list of quantity-payment pairs $\{q_{jv}, z_{jv}\}_{v \in V}$. It is not restrictive to assume that the number of items in this list is equal to the number of ex post buyer types. A seller's strategy is summarized by a price schedule $Z_j : Q_j \to \Re$ where $Q_j = \{q_{jv}\}_{v \in V}$ and $z_{jv} = Z_j(q_{jv})$ for all $v \in V$. See Maskin and Riley (1984) which analyzes the problem of a monopolistic seller who faces no price competition. In their model, however, there are no matching frictions.

¹¹Buyers can also play mixed strategies and randomize over price schedules for which they are indifferent.

¹²By the revelation principle, the formulation a seller announces a price schedule and buyers with different realizations of ε self-select along this schedule by choosing their most preferred price quantity combination is equivalent to a formulation where the seller announces a direct revelation mechanism which induces truthful revelation by the buyers. See Section 4.

 $^{^{13}}$ See Pissarides (1990).

where \mathcal{M} is a standard matching function that maps the measures of buyers and sellers in submarket j onto the measure of trading meetings in this submarket.

An interesting special case of m^s is:

$$m^{s}(\theta_{j}) = 1 - \exp(-\theta_{j}). \tag{3}$$

This case arises if buyers use identical mixed strategies to select a seller among those who post equivalent price offers and, because selling is time consuming, each seller can serve at most one customer (see Peters (2000)). This is the typical setup in the frictional assignment literature. Another interesting special case is

$$m^{s}(\theta_{j}) = \frac{\theta_{j}}{1 + \theta_{j}}.$$
(4)

This special case arises if each buyer is randomly matched with a trader (buyer or seller) in the submarket where search is directed. (In this case, $m^s(\theta_j)$ is equal to the fraction of buyers over traders in submarket j).

We consider an environment where buyers and sellers are ex-ante symmetric. Given that buyers are free to choose among different price schedules, in a directed search equilibrium, all buyers must attain the same expected payoff. Furthermore, in equilibrium no seller should have an incentive to deviate by posting a different price schedule. To attract buyers, the offer of the deviating seller must yield at least the common expected payoff buyers attain in equilibrium.¹⁴ Therefore, in a directed search equilibrium the price schedules and the degree of congestion associated with them must maximize the expected payoff of sellers subject to the constraint that the buyers get the common expected payoff. As we show in Section 4, this implies that in our framework there is a single price schedule and a single submarket in equilibrium.

We proceed to develop our model with the following steps. In Sections 3, 4, and 5, we construct a static version of the model. In Section 3, we analyze the optimal behavior of a household for a given price schedule Z(q). Section 4 studies the endogenous determination of the price schedule Z(q) under directed search. Section 5 collects the results of Sections 3 and 4 in a general equilibrium model where both the behavior of the household and prices are endogenous.

¹⁴Since buyers can play mixed strategies, the expected number of buyers attracted by a single deviator need not be an integer.

3 A Representative Household

In this section, we describe the behavior of a household whose buyers and sellers trade in a retail market with a given price schedule Z(q) and a given degree of congestion θ . Because there is a continuum of households in the economy, the household takes the price schedule and the degree of congestion as given. We assume Z(q) to be continuously differentiable and concave with Z(0) = 0. The next section endogeneizes Z(q) and validates these assumptions. We adopt the following notation. Lower-case letters denote the decision variables of the household. Upper-case letters denote the decisions of the other households and hence aggregate quantities, which are also taken as given by our household. In a symmetric equilibrium, lower-case letters are equal to the corresponding upper-case letters.

The household is composed of an infinite number of individuals. The individuals of the household are assigned to one of three different tasks: production (producers), purchase (buyers), and sale of commodities (sellers). Each individual is endowed with one unit of labor and can perform at most one task. The timing of the model is as follows. First, the members of the household are divided into producers, buyers, and sellers. The measure of individuals assigned to each activity is denoted by n, b, and s, respectively. Second, all members perform their assigned activities. Specifically, producers use their labor n to generate output, which immediately becomes available for sale. Buyers visit sellers from other households. Upon meeting a seller, they experience a random shock ε that scales the incremental utility that the good brings to the household. Shocks are uniformly distributed in [0, 1] and are independent across trading meetings. While the distribution of these shocks in public information, the realized value of ε is the buyer's own private information. Contingent on the realized value of ε , buyers choose the quantity q_{ε} they want to acquire from the seller, and pay $z_{\varepsilon} = Z(q_{\varepsilon})$.¹⁵ This payment is immediately debited from the household's account in the clearing-house. Sellers go to retail outlets where they wait for buyers to visit them. When a seller of the household is visited by a buyer who purchases a quantity Q_{ε} , the seller fetches this quantity from the current household production and collects a payment $Z_{\varepsilon} = Z(Q_{\varepsilon})$, which is immediately transferred to the household's account in the clearing-house. Finally, once all their tasks are completed, all the individuals of the household get together and equally share the consumption of the goods purchased.

The assumption that there is an infinite number of members in the household ensures full

¹⁵We use ε subscript to denote that q and z are functions of ε .

diversification of trading risks. Further, we assume that this number is countable to ensure that the set of buyers in the household is of measure zero relative to the set of goods for sale by other households.¹⁶ This assumption is technically important because it implies that there are no repeated purchases of the same good with probability one.¹⁷ For convenience, we normalize the measure of individuals in the household to one, that is we measure sets of individuals in the household as fractions of the individuals in the set relative to the individuals in the household.

The household's utility is $\mathcal{U}(c)$, where $\mathcal{U} : \Re_+ \to \Re$ is continuously differentiable, strictly increasing, and concave. The variable c is a hedonic measure of consumption, which satisfies

$$c = \left[bm^{b}(\theta) \int_{0}^{1} \varepsilon q_{\varepsilon}^{1-\sigma} d\varepsilon \right]^{\frac{1}{1-\sigma}}, \quad \sigma \in (0,1).$$
(5)

The term $bm^b(\theta)$ is the measure of goods located by the *b* buyers in the household who on average visit $m^b(\theta)$ sellers each. The law of large numbers implies that, ex post, the realizations of the preference shocks for these goods are uniformly distributed in [0, 1]. This consumption aggregator implies that *c* doubles when q_{ε} doubles for all ε , and *c* more than doubles when *b* doubles. Therefore, the household has a preference for variety,¹⁸ which increases with the parameter σ and vanishes when $\sigma \to 0$. The Appendix derives the hedonic measure of consumption in (5) from a standard Dixit-Stiglitz aggregator. Finally, to avoid dealing with uninteresting corner solutions, we assume that \mathcal{U} is sufficiently concave, so utility is a concave function of *b*.

Production y depends on the amount of labor n employed:

$$y = f(n). \tag{6}$$

The production function $f : \Re_+ \to \Re_+$ is assumed continuously differentiable, strictly increasing, and concave. The sales of the household are equal to the measure of buyers contacted by the sellers of the household times the average quantity sold by each seller, $\overline{Q} = \int_0^1 Q_{\varepsilon} d\varepsilon$. Since there is no aggregate uncertainty in a large household, the amount of output both produced and sold is identical:

$$y = sm^s(\theta)\overline{Q}.\tag{7}$$

¹⁶Remember that there is a continuum of households and thus a continuum of differentiated goods that each household can purchase.

¹⁷Also, the standard Law of Large Numbers assumes a countable infinite number of random variables.

¹⁸Note that the household does not buy infinitesimal amounts of the goods it purchases because it is costly to send extra buyers to the market.

The household must satisfy the budget constraint, that is, total expenditure cannot exceed total sales revenue:

$$sm^s(\theta)\overline{Z} - bm^b(\theta)\overline{z} \ge 0.$$
 (8)

Here, $\overline{Z} = \int_0^1 Z(Q_{\varepsilon}) d\varepsilon$ is the average revenue of a seller, and $\overline{z} = \int_0^1 Z(q_{\varepsilon}) d\varepsilon$ is the average expenditure of a buyer.

The measure of individuals assigned to the three different activities in the household must add up to one:

$$b + s + n = 1. \tag{9}$$

The household chooses $\{q_{\varepsilon}\}_{\varepsilon \in [0,1]}$, b, s, n, c and y to maximize \mathcal{U} subject to (5) to (9), and non-negativity constraints for all variables. Condition (7) can be substituted into (6) to form a single resource constraint. Also, (5) can be substituted into the objective function $\mathcal{U}(c)$. Using Lagrange multipliers μ , λ , and ϑ for the resource constraint, the budget constraint, and the labor allocation constraint respectively, the first-order conditions for an interior maximum are:

$$\mathcal{U}'(c)c^{\sigma}\varepsilon q_{\varepsilon}^{-\sigma} = \lambda \mathcal{Z}'(q_{\varepsilon}) \text{ for } \varepsilon \in [0,1],$$
(10)

$$m^{b}(\theta) \int_{0}^{1} \left[\frac{\varepsilon \mathcal{U}'(c) c^{\sigma} q_{\varepsilon}^{1-\sigma}}{1-\sigma} - \lambda z_{\varepsilon} \right] d\varepsilon = \vartheta,$$
(11)

$$m^{s}(\theta) \int_{0}^{1} \left(\lambda Z_{\varepsilon} - \mu Q_{\varepsilon}\right) d\varepsilon = \vartheta, \text{ and}$$
 (12)

$$\mu f'(n) = \vartheta. \tag{13}$$

Condition (10) states that, conditional on each possible realization ε , buyers must equate the marginal utility of purchasing an extra unit of a good with the marginal value of the payment required in return. The other conditions imply that the value of the marginal product of labor in all three occupations must be the same. This common value is ϑ . Condition (11) equates ϑ to the expected consumer surplus obtained by a buyer. Condition (12) equates ϑ to the expected surplus generated by a seller. Finally, condition (13) equates ϑ to the value of the marginal product of labor of a producer.

4 Commercial Pricing

In this section, we model the interaction between buyers and sellers in the marketplace in a directed search equilibrium when sellers do not know how much buyers are willing to pay for their merchandise. A bilateral trade is described by a pair $(q, z) \in \Re^2_+$ specifying the quantity q supplied by the seller and the payment z given in return by the buyer. In the previous section, the two elements of this vector were linked by an exogenous price schedule $z = \mathcal{Z}(q)$. The purpose of this section is to endogeneize \mathcal{Z} .

The revelation principle allows us to reformulate the problem where sellers post price schedules and buyers, with different realizations of the preference shock, self-select along the schedule of the sellers they are matched with. Without loss of generality, we may assume that sellers posts direct revelation mechanisms that induce truth-telling by the buyers (that is, incentive compatible mechanisms). Feasible mechanisms must also satisfy the individual rationality constraints of the buyers. Buyers observe the mechanisms posted by the sellers and decide to trade according to a particular mechanism. In a directed search equilibrium, all posted mechanisms and the degrees of congestion associated with them must maximize the expected payoff of the seller subject to the constraint that buyers get a common expected payoff. This result has two direct implications. First, the mechanisms posted by sellers must be optimal in the set of incentive compatible and individually rational mechanisms. That is, any posted mechanism must be incentive efficient. Second, the congestion associated with the mechanism must be such that no other combination of an incentive efficient mechanism and a congestion level yield the same expected payoff to the buyer and higher expected payoff to the seller.

4.1 Incentive Efficient Direct Revelation Mechanisms

We begin by characterizing the set of incentive efficient direct revelation mechanisms. That is, we characterize the direct mechanisms that maximize a weighted sum of the expected payoffs of the buyer and the seller subject to the incentive compatibility and individual rationality constraints. We show that the outcome of each incentive efficient direct revelation mechanisms is implemented when sellers post a price schedule in a particular class, and buyers who choose to trade at these prices select their most preferred price-quantity combination depending on the realization of the preference shock.¹⁹

In a meeting between a buyer and a seller, the buyer knows the realized value of ε , but the seller does not. All the seller knows is that ε is a random variable uniformly distributed on the interval [0, 1]. We refer to the realization of ε as the buyer's *type*.

All traders maximize the objectives of the households they belong to. Specifically, the incre-

¹⁹One element in this class corresponds with the monopolistic price schedule derived in Maskin and Riley (1984) and Mussa and Rosen (1978).

mental utility of a buyer's purchase to the household is the consumer surplus:

$$\mathcal{U}^{b}(q,z;\varepsilon) = \frac{\varepsilon \psi q^{1-\sigma}}{1-\sigma} - \lambda z, \qquad (14)$$

where

$$\psi \equiv \mathcal{U}'(c)c^{\sigma},\tag{15}$$

and λ is the Lagrange multiplier associated with (8). The incremental utility of a seller's sale is the gross commercial margin:

$$\mathcal{U}^{s}\left(q,z\right) = \lambda z - \mu q,\tag{16}$$

where μ is the multiplier associated with (7).²⁰ By symmetry, the variables ψ , λ , and μ are identical across households. Because each buyer and each seller are infinitesimal in the household they belong to, the values of ψ , λ , and μ are not affected by their individual actions. For ease of exposition, we shall refer to \mathcal{U}^b and \mathcal{U}^s as the utility functions of buyers and sellers respectively.

A direct revelation mechanism is a schedule of type-contingent trades $\{q_{\varepsilon}, z_{\varepsilon}\}_{\varepsilon \in [0,1]}$. A direct revelation mechanism is incentive efficient if it maximizes a weighted sum of the expected utilities of buyers and sellers:

$$\{q_{\varepsilon}, z_{\varepsilon}\}_{\varepsilon \in [0,1]} = \arg \max \left[(1-\omega) \int_0^1 \mathcal{U}^b \left(q_{\varepsilon}, z_{\varepsilon}; \varepsilon\right) d\varepsilon + \omega \int_0^1 \mathcal{U}^s \left(q_{\varepsilon}, z_{\varepsilon}\right) d\varepsilon \right]$$
(17)

where $\omega \in [0, 1]$, subject to the following constraints:

1. Incentive Compatibility: Buyers must have no incentive to lie about their type:

$$\varepsilon' \in \arg \max_{\varepsilon \in [0,1]} \left[\mathcal{U}^b \left(q_{\varepsilon}, z_{\varepsilon}; \varepsilon' \right) \right], \text{ for all } \varepsilon' \in [0,1].$$
 (18)

2. Individual Rationality: Buyers and sellers must receive non-negative utility in all meetings:²¹

$$\mathcal{U}^{b}(q_{\varepsilon}, z_{\varepsilon}; \varepsilon) \ge 0, \text{ for all } \varepsilon \in [0, 1], \text{ and}$$
 (19)

$$\mathcal{U}^{s}(q_{\varepsilon}, z_{\varepsilon}) \ge 0, \text{ for all } \varepsilon \in [0, 1].$$
 (20)

In the appendix, we use standard arguments in the mechanism design literature to characterize the set of incentive efficient direct revelation mechanisms. This solution is summarized in the following proposition:

 $^{^{20}\}mathcal{U}^b$ and \mathcal{U}^s are defined so that they are aligned with the household's first order conditions (11) and (12).

²¹This property follows from the combination of two facts: the buyer observes the variety offered by the seller as soon as they meet, and both buyers and sellers have the option of not trading. If sellers could hide their variety for sale, then they could demand a payment from the buyers to reveal this information. In this case, buyers could end up with negative utility in some meetings.

Proposition 1 The incentive efficient direct revelation mechanisms which solve program (17) subject to (18) to (20), are the following:

$$q_{\varepsilon} = \begin{cases} 0 & \text{for } \varepsilon \in [0, \gamma) \\ \left(\frac{\psi}{\mu} \frac{\varepsilon - \gamma}{1 - \gamma}\right)^{\frac{1}{\sigma}} & \text{for } \varepsilon \in [\gamma, 1] \end{cases}$$
(21)

$$z_{\varepsilon} = \frac{1}{\lambda} \left[\gamma \frac{\psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} + (1-\gamma) \, \mu q_{\varepsilon} \right], \qquad (22)$$

where

$$\gamma = \begin{cases} 0 & \text{for } \omega \in [0, 0.5) \\ \frac{2\omega - 1}{3\omega - 1} & \text{for } \omega \in [0.5, 1] \end{cases}$$
(23)

The following proposition establishes that the outcome of an incentive efficient direct revelation mechanism can be implemented if sellers post a simple non-linear price schedule and buyers optimally choose the quantity they want to acquire at these prices.

Proposition 2 The outcome of an incentive efficient direct revelation mechanism can be attained when buyers choose q facing the following price schedule:

$$\mathcal{Z}(q) = \frac{1}{\lambda} \left[\gamma \frac{\psi q^{1-\sigma}}{1-\sigma} + (1-\gamma) \, \mu q \right], \tag{24}$$

where $\gamma \in [0, 0.5]$ satisfies (23). We refer to a price schedule with parameter γ as the price schedule γ .

As long as commercial margins are positive ($\gamma > 0$), the pricing schedule $\mathcal{Z}(q)$ is strictly concave. This concavity implies that the per unit price of goods declines with q, or equivalently buyers obtain quantity discounts. This is not unrealistic. In retail trade, sometimes we observe explicit quantity discounts, but most often quantity discounts are implicit in the packaging of products (the larger is the box of nails, the lower is the per unit price).

When buyers have full market power, that is when $\omega = 0$, prices cover only the marginal cost of production ($\gamma = 0$). In this case, buyers capture the whole trading surplus and the individual rationality constraint of the seller binds. This constraint continues to bind as long as $\omega \leq 0.5$. For $\omega > 0.5$, commercial margins are positive and both the buyer and the seller appropriate a fraction of the trading surplus. Even when sellers have full market power, that is when $\omega = 1$ and $\gamma = 0.5$, they are not able to extract the whole trading surplus because they do not know their clients' type. The latter case is an interesting benchmark, which we denote monopolistic search. With monopolistic search, sellers post the prices that maximize their expected profits in an environment where their price schedules have no effect on the number of clients visiting their outlets. This might be the relevant equilibrium concept for some tourist areas where buyers have little knowledge about where to shop. However, in most commercial areas sellers are aware that by posting low prices they can attract clients to their outlets. To capture these competitive pressures on prices, we adopt directed search as our concept of equilibrium.

4.2 Trade-off between prices and congestion

Let $\mathcal{V}^{b}(\theta_{j}, \gamma_{j})$ and $\mathcal{V}^{s}(\theta_{j}, \gamma_{j})$ respectively be the expected utility of buyers and sellers in a submarket with congestion θ_{j} and price schedule γ_{j} :

$$\mathcal{V}^{b}(\theta_{j},\gamma_{j}) = m^{b}(\theta_{j}) \int_{0}^{1} \left(\frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} - \lambda z_{\varepsilon}\right) d\varepsilon, \text{ and}$$
(25)

$$\mathcal{V}^{s}(\theta_{j},\gamma_{j}) = m^{s}(\theta_{j}) \int_{0}^{1} \left(\lambda z_{\varepsilon} - \mu q_{\varepsilon}\right) d\varepsilon$$
(26)

where q_{ε} and z_{ε} are given by (21) and (22). Integrating these expressions we obtain

$$\mathcal{V}^{b}(\theta_{j},\gamma_{j}) = m^{b}(\theta_{j}) \frac{\sigma^{2}}{1-\sigma^{2}} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}} (1-\gamma_{j})^{2}, \text{ and}$$
(27)

$$\mathcal{V}^{s}(\theta_{j},\gamma_{j}) = m^{s}(\theta_{j}) \frac{2\sigma^{2}}{1-\sigma^{2}} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}} \gamma_{j}(1-\gamma_{j}).$$

$$(28)$$

In a directed search equilibrium, \mathcal{V}^b must be common to all active submarkets, because buyers are free to choose the submarket where they direct their search. Also, the pair (θ_j, γ_j) must maximize $\mathcal{V}^s(\theta_j, \gamma_j)$ subject to the constraint that $\mathcal{V}^b(\theta_j, \gamma_j)$ is equal to the common expected utility attained by buyers in all active submarkets. The solution to this maximization problem, if one exists, is unique.²² Hence, there is a unique submarket in a directed search equilibrium, and the subscript *j* can be omitted. As long as sellers capture a fraction of the expected trading surplus,²³ the equilibrium pair (θ, γ) is interior and satisfies the following first order condition:

$$\frac{\frac{\partial \mathcal{V}^{b}(\theta,\gamma)}{\partial \theta}}{\frac{\partial \mathcal{V}^{b}(\theta,\gamma)}{\partial \gamma}} = \frac{\frac{\partial \mathcal{V}^{s}(\theta,\gamma)}{\partial \theta}}{\frac{\partial \mathcal{V}^{s}(\theta,\gamma)}{\partial \gamma}}.$$
(29)

²²Substituting the constraint in the objective function and dropping unnecessary constants, θ_j must maximize the strictly concave function: $m^s(\theta_j) \left\{ 1 - \left[vm^b(\theta_j) \right]^{-0.5} \right\} \left[vm^b(\theta_j) \right]^{-0.5}$, where v is a constant inversely proportional to the given value for \mathcal{V}^b .

 $^{^{23}}$ In our model, a zero expected payoff for sellers cannot be an equilibrium because there would be no sellers. See next section.

That is, buyers and sellers must have a common marginal rate of substitution of θ for γ . Differentiating (27) and (28) and substituting into (29), we obtain

$$\gamma = \frac{1 - \eta(\theta)}{2},\tag{30}$$

where $\eta(\theta)$ is the elasticity of the function m^s :

$$\eta(\theta) = \frac{m^{s'}(\theta)\theta}{m^s(\theta)}.$$
(31)

Equation (30) provides a simple restriction that pricing schedules and congestion levels must satisfy in a directed search equilibrium. In the next section, we pinpoint the equilibrium pair (θ, γ) by combining this restriction with the optimal choice of the measure of buyers and sellers in each household.

5 General Equilibrium

In this section, we characterize a directed search equilibrium by combining the optimal behavior of households in Section 2 with the endogenous determination of the pricing function in Section 3. **Definition:** The tuple $\{\theta, n, b, s, c, y, q_{\varepsilon}, z_{\varepsilon}, \psi, \mu, \lambda, \mathcal{Z}(q)\}$ is a symmetric directed search equilibrium if

1. All households choose $\{n, b, s, c, y, q_{\varepsilon}\}$ taking as given the payment schedule $\mathcal{Z}(q)$ and the market congestion ratio θ , and the implied payments are $z_{\varepsilon} = \mathcal{Z}(q_{\varepsilon})$.

2. The payment schedule $\mathcal{Z}(q)$ satisfies (24) when ψ satisfies (15) and μ and λ are the Lagrange multipliers associated with the resource constraint and the budget constraint in the household's optimization program.

- 3. The pair (γ, θ) satisfies condition (30).
- 4. Congestion ratios are consistent with individual behavior: $\theta = b/s$.

Since in a symmetric search equilibrium all households behave identically, lower-case and upper-case letters coincide.

In equilibrium, a household must obtain the same utility from allocating an individual to any of the three activities, as it is implied by the first-order conditions (11) to (13). Therefore, $\mathcal{V}^b(\theta, \gamma)$ and $\mathcal{V}^s(\theta, \gamma)$ must be equal. Using (27) and (28) together with (2), this equality implies

$$\theta = \frac{1 - \gamma}{2\gamma}.\tag{32}$$

According to (32), the congestion of buyers in the market θ is inversely related to γ . As γ increases buyers capture a smaller fraction of the trade surplus (see Proposition 2), so households respond by sending fewer buyers and more sellers to the marketplace. Combining (32) with (30), we obtain that congestion θ in a directed search equilibrium is the solution to

$$\frac{1}{\frac{1}{2}+\theta} = 1 - \eta(\theta). \tag{33}$$

Consequently, the efficiency condition (29) together with the equality of expected payoffs of buyers and sellers determine the two key variables in the retail sector: congestion θ and price schedule γ .

To determine the remaining variables in the model, we must combine the resource constraints, the household's first order conditions, and the definitions of θ , c, and ψ . Equating $\mathcal{V}^b(\theta, \gamma)$ to the marginal payoff of a producer, $\mu f'(n)$, we obtain

$$f'(n) = m^{s}(\theta) \frac{2\sigma^{2}}{1 - \sigma^{2}} \left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}} \gamma \left(1 - \gamma\right).$$
(34)

Integrating the quantities traded when buyers face a price schedule γ which are given by (21), we obtain the average sales

$$\overline{q} = \frac{\sigma}{1+\sigma} \left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}} (1-\gamma).$$
(35)

Combining (6), (7), (34), and (35) gives

$$\frac{f'(n)}{f(n)} = \frac{2\sigma}{1-\sigma}\frac{\gamma}{s}.$$
(36)

Using the equilibrium condition $\theta = b/s$, and the labor resource constraint (9), we have

$$s = \frac{1-n}{1+\theta}.\tag{37}$$

Finally, combining (36) with (37) and using (32), we obtain

$$\frac{(1-n)f'(n)}{f(n)} = \frac{\sigma(1+\gamma)}{1-\sigma}.$$
(38)

Equation (38) determines n for a given value of γ . In a directed search equilibrium, γ and θ are obtained from (30) and (33). Given θ , γ , and n, the equilibrium values for s, b, ψ/μ , q_{ε} , c, y, ψ , and μ are recursively determined by (36), (1), (34), (21), (5), (6), (15), and (34). The utility value of the payment schedule $\lambda \mathcal{Z}(q)$ is determined by (24), but the precise values of λ and $\mathcal{Z}(q)$ are indeterminate because they depend on the units in which payments are measured.

The following proposition summarizes the conditions for the existence and uniqueness of a directed search equilibrium (the proof is in the Appendix):

Proposition 3 If the elasticity $\eta(\theta)$ is non-increasing and belongs to the interval (0,1),²⁴ and the following terminal conditions are satisfied: f(0) = 0 and f'(0) > 0, then a directed search equilibrium exists and is unique.²⁵

To calculate the average commercial margin in a directed search equilibrium, we use (21) to obtain the following relationship:

$$\psi \int_0^1 \varepsilon q_{\varepsilon}^{1-\sigma} d\varepsilon = (1+\gamma\sigma)\,\mu\overline{q}.\tag{39}$$

Substituting (39) and (32) into (11) and (12), we get

$$\frac{2\gamma}{1-\gamma} \left(\frac{1+\gamma\sigma}{1-\sigma} \mu \overline{q} - \lambda \overline{z} \right) = \lambda \overline{z} - \mu \overline{q}.$$
(40)

Using (40), we obtain that the average commercial margin

$$\frac{\lambda \overline{z} - \mu \overline{q}}{\mu \overline{q}} = \frac{\sigma}{1 - \sigma} \left[1 - \eta \left(\theta \right) \right]. \tag{41}$$

The commercial margin increases with both the preference for diversity σ and the contribution of sellers in the matching process measured by $1 - \eta (\theta)$.²⁶

5.1 Multiple Matches

In our model, there is no logical reason to restrict buyers and sellers to at most one match each period. The extension to multiple matches is straightforward for the following reasons. Individual traders (buyers or sellers) are infinitesimal in the household where they belong, so their individual actions have no effect on the shadow valuations of goods, labor, and payments in the household: ψ , μ , and λ . Moreover, thanks to the clearing-payments mechanism, an individual buyer is not restricted by the money he or she carries. Likewise, thanks to the immediate access to current production, an individual seller is not restricted by available commercial inventories. Therefore, the optimal strategies of buyers and sellers in successive trading meetings do not depend on their

 $^{^{24}}$ The two examples (3) and (4) satisfy these conditions.

²⁵This proof also uses the continuity and concavity of m^s and f, and the continuity and convexity of m^b assumed throughout the paper.

²⁶With full information, $1 - \eta(\theta)$ is the fraction of the trade surplus appropriated by sellers in a directed search equilibrium. This division of the surplus endogeneizes the matching externalities, so the equilibrium allocation is efficient. See Hosios (1990).

individual histories. Our analysis and results apply with almost no modification to a generalized model where buyers and sellers can perform multiple matches each period. In that model, we reinterpret $m^s(\theta)$ and $m^b(\theta)$ as the expected number of trading matches that a seller and a buyer perform during a period respectively. With this reinterpretation, the image of these two functions is $[0, \infty)$ instead of [0, 1].²⁷ Conveniently, this allows for a Cobb-Douglas matching technology for which the elasticity of $m^s(\theta)$ is a constant.

6 Welfare

This section studies the welfare properties of the directed search equilibrium. To this end, it characterizes the optimal allocation that a benevolent central planner would choose in order to maximize the utility of a representative household. Following standard practice, the central planner is not only bound by the resources available in the economy, but also by the bilateral matching among traders. In addition, it is sensible that we restrict the central planner to information that is publicly available. However, imposing this restriction is much more subtle than it first appears. The incentive compatibility constraint (18) arising from the fact that the buyers' types are private information is only binding when selling costs must be financed with the revenue from sales (see below). Moreover, the central planner can easily affect this revenue with policy tools that, in principle, are respectful to private information, for example, a sales tax or a sales subsidy. Due to the subtleties of restricting the central planner to public information, we start by characterizing the first best allocation, in which the central planner is not bound by the incentive compatibility constraints imposed by private information. This allocation is an interesting benchmark in itself and is useful to evaluate the welfare costs of private information. Later, we introduce private information with specific assumptions about the policy tools at the disposal of the central planner.

6.1 The Optimum with Complete Information

The first best allocation is one in which the central planner maximizes the utility of a representative household subject to the resources available in the economy and the bilateral matching among traders. Let $\mathcal{M}(b,s)$ be the aggregate matching function, that is $\mathcal{M}(b,s) = bm^b(b/s) = sm^s(b/s)$.

²⁷With multiple matches, it is important that buyers can play mixed strategies when they choose which prices to search. Otherwise, a deviating seller would have a hard time attracting prospective buyers because these would dislike being in a submarket where there is a single good.

The planner must solve the following program:

$$\max_{b,s,n,q_{\varepsilon}} \mathcal{U}(c), \text{ where } c = \left[\mathcal{M}(b,s) \int_{0}^{1} \varepsilon q_{\varepsilon}^{1-\sigma} d\varepsilon\right]^{\frac{1}{1-\sigma}},$$
(42)

where subject to (9) and

$$\mathcal{M}(b,s) \int_{0}^{1} q_{\varepsilon} d\varepsilon = f(n) \,. \tag{43}$$

The first-order conditions of this problem are:

$$\varepsilon \psi q_{\varepsilon}^{-\sigma} = \mu, \tag{44}$$

$$\mathcal{M}_b(b,s) \int_0^1 \frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} d\varepsilon - \mu \mathcal{M}_b(b,s) \int_0^1 q_{\varepsilon} d\varepsilon = \mu f'(n), \tag{45}$$

$$\mathcal{M}_{s}(b,s)\int_{0}^{1}\frac{\varepsilon\psi q_{\varepsilon}^{1-\sigma}}{1-\sigma}d\varepsilon - \mu\mathcal{M}_{s}(b,s)\int_{0}^{1}q_{\varepsilon}d\varepsilon = \mu f'(n), \tag{46}$$

where μ is the Lagrange multiplier of (43) and $\psi = \mathcal{U}'(c)c^{\sigma}$. In the first best, the marginal utility of consuming each good must be equal to the marginal production cost. Also, the marginal social benefit of employment in all three activities must be the same. Comparing (45) and (46), we obtain

$$\mathcal{M}_b(b,s) = \mathcal{M}_s(b,s). \tag{47}$$

Equality (47) implies that for a given number of traders the number of trading meetings is maximized. Equations (9) and (43) to (46) can be easily manipulated to obtain an almost explicit solution of the first best allocation (see the Appendix for details):

Proposition 4 The first best allocation, in which the central planner has both complete information and control over all variables, is characterized by the following equations:

$$\frac{1}{1+\theta} = 1 - \eta(\theta), \tag{48}$$

$$\frac{(1-n)f'(n)}{f(n)} = \frac{\sigma}{1-\sigma},$$
(49)

$$s = \frac{1-n}{1+\theta},\tag{50}$$

$$b = \theta s, and$$
 (51)

$$q_{\varepsilon} = \frac{1+\sigma}{\sigma} \frac{f(n)}{\mathcal{M}(b,s)} \varepsilon^{\frac{1}{\sigma}}.$$
(52)

The following proposition compares the allocations in the first best and in the directed search equilibrium. To facilitate this comparison, it specializes the production and matching technologies to standard functional forms: **Proposition 5** Suppose all the assumptions used in Proposition 3 to ensure existence and uniqueness of an equilibrium hold. Let an asterisk denote first best and no asterisk denote directed search equilibrium. The following relations hold:

$$\frac{b}{s} > \frac{b^*}{s^*},\tag{53}$$

$$n < n^*. (54)$$

In addition, if the f is isoelastic,

$$\frac{s}{n} \ge \frac{s^*}{n^*}$$
, with equality if $\eta(\theta)$ is constant. (55)

Finally, if f is constant returns to scale,

$$\overline{q} < \overline{q}^*. \tag{56}$$

In a directed search equilibrium, selling costs are financed with commercial margins that create a wedge between the marginal production cost of merchandises and the price paid by buyers. This wedge reduces the quantities that buyers purchase in each transaction for given valuations of goods and labor ψ and μ (compare (21) and (44)). However, ex-ante price competition among sellers narrows commercial margins so the congestion of buyers to sellers in equilibrium is higher than in the first best. Intuitively, the equilibrium price schedule must play two conflicting allocational roles: It must signal buyers the opportunity cost of the goods they are considering to purchase, and it must finance retail costs in order to induce an efficient ratio of sellers over buyers. The equilibrium price schedule settles on a compromise between these two roles. Prices are higher than the social opportunity cost of goods but not high enough to finance the first best ratio of sellers over buyers. Even though, commercial margins induce buyers to purchase smaller quantities, the comparison between q_{ε} and q_{ε}^* is ambiguous because the equilibrium and the first best assign a different shadow value to the cost of production relative to the utility of consumption (μ/ψ) . In the special case where f is constant returns to scale, we can prove that on average buyers purchase an inefficiently low quantity when they meet a seller.

In the baseline case where both f and m^s are isoelastic, the ratio (s/n) is identical in the equilibrium and the first best allocations. Therefore, in this baseline case the inefficiently high (b/s) ratio implies that in equilibrium the measure of buyers is inefficiently high and the measures of sellers and producers are inefficiently low. In general, the equilibrium allocation of labor relative to the first best depends on whether the elasticities f and m^s are increasing or decreasing with respect to n and θ respectively.

As in Moen (1997) and Shimer (1996), the first best allocation can be implemented as a directed search equilibrium when buyers and sellers have complete information. In this case, a bilateral trade must maximize the joint surplus of buyers and sellers. Moreover, the ex-ante price competition among sellers leads to sharing this surplus according to Hosios's (1990) rule, that is buyers get a fraction $\eta(\theta)$ of the surplus and sellers get $1 - \eta(\theta)$, and θ is the first best level of congestion. With private information about buyers' types, this equilibrium breaks down because the incentive compatibility constraint (18) is violated.

The first best allocation can be decentralized as a directed search equilibrium where buyers' types are private information if sellers can charge a flat fee to prospective buyers prior to the realization of the preference shock ε . In this case, the fee covers selling costs without having to add a positive commercial margin on the price of merchandises. In our model, such a fee is prevented by the assumption that buyers get the preference shock as soon as they meet a seller, so a buyer is not willing to pay the fee if the realization of ε is low. That is, flat fees violate the individual rationality constraint of buyers. More generally, one could realistically assume that the buyers' satisfaction from a commercial transaction depends on the service effort provided by the seller. In this more complicated model, the flat unconditional fee is also discouraged by the moral hazard problem it generates on the effort exercised by sellers. In reality, we find flat fees in warehouse clubs. However, sales in warehouse clubs are a small fraction of the economy-wide retail sales, and even these clubs charge fees that cover a small fraction of their commercial costs.

The following proposition summarizes these two ways of decentralizing the first best allocation:

Proposition 6 The first best allocation where the planner has complete information can be implemented as a directed search equilibrium if buyers and sellers have complete information or if sellers can charge a lump-sum fee to prospective buyers prior to the realization of the preference shock ε .

6.2 The Optimum with Private Information

This subsection characterizes second best allocations when the central planner has limited information. Given that buyers' types are private information it is natural to assume that the central planner cannot directly observe them. However, if the central planner can monitor the allocation of labor in each household, then the first best allocation can be easily implemented by dictating to households the allocation of labor, charging buyers the marginal cost of producing merchandise, and transferring the proceeds of these payments to sellers. With this mechanism, buyers truthfully reveal their types without any efficiency loss. To make our analysis more realistic and more interesting than this simple result, we assume that the central planner cannot directly observe how households allocate their labor. This assumption is in line with the unobservability of leisure (and so the allocation of time when there are two activities) in the standard theory of taxation.

Specifically, we assume that the planner can only observe market transactions and only has control of the direct revelation mechanisms by which transactions are conducted, or equivalently the price schedules faced by buyers and sellers. Given these price schedules, buyers choose the quantities they want to purchase, which sellers supply as long as prices are above the marginal costs of production. Households decide the allocation of labor taking into account the expected returns from each activity.

In principle, buyers and sellers could face different price schedules. If so, the gap between the two schedules is a sales tax (or a sales subsidy) to be collected (or distributed) by a government who balances its budget with a lump-sum subsidy (or a lump-sum tax) on households. In the two propositions that follow, we make alternative assumptions on the central planner's ability to control the pricing mechanism and to impose lump-sum taxes on households.

In the absence of any restriction on these policy instruments, the central planner can implement the first best allocation despite the presence of private information:

Proposition 7 If the central planner can resort to lump-sum taxes on households to finance a linear sales subsidy, then the first best allocation characterized in Proposition 4 can be decentralized with the following price schedule faced by buyers:

$$\mathcal{Z}(q) = \frac{\eta(\theta)\mu q}{\lambda},\tag{57}$$

and a linear sales subsidy at the gross rate:

$$T = \frac{1 - \sigma \eta(\theta)}{(1 - \sigma) \eta(\theta)},\tag{58}$$

where θ is the solution to (48). (The price schedule faced by sellers is TZ(q)).

This proposition implies that the planner could not do better by introducing additional control instruments. Intuitively, the proposition holds because the planner picks the tab for the selling costs with the sales subsidy. Hence, the prices faced by the buyers can be equated to the marginal social cost of production of merchandises. The implementation of the first best allocation in Proposition 7 depends on the existence of lump-sum taxes. In the model, this is not problematic because all households are identical. However, this homogeneity is unrealistic and has been assumed only for simplicity. The next proposition assumes the absence of lump-sum taxes.

Proposition 8 If the only policy instrument is to set the retail trade price schedule and the marginal product of labor is constant (the production function is affine), a directed search equilibrium is optimal.

If the production function is affine,²⁸ a central planner without recourse to lump-sum taxes cannot improve upon the directed search equilibrium. When the production function is not affine, regulation of the price schedule can be welfare enhancing because it has an indirect effect on the marginal product of labor which differs from its first best value. However, as long as the production function is differentiable, it can always be approximated with an affine production function, and so the welfare gains of deviating from the directed search equilibrium are second order of magnitude. In the next section, we use numerical simulations to check that these gains are very small.

7 Commerce in a Neoclassical Growth Framework

This section sketches how to embed the model developed in the previous sections in a Neoclassical growth framework. In the extended model, time is infinite and discrete. Each period in the life of a household proceeds as in the static model. Also, in the resulting synthesis, the economy has two sectors: one produces goods combining capital and labor as in the Neoclassical growth model, the other exchanges goods in directed search markets where buyers' types are private information.²⁹ To make our model as close as possible to the basic Neoclassical growth model, we assume that investment does not require installation or commercial costs. However, future work could incorporate these features. For brevity, we omit all proofs, which are either standard in the Neoclassical growth theory or parallel the arguments in the previous sections.

The household's intertemporal utility function is

 $^{^{28}}$ Since labor is the only production input, this assumption is satisfied with constant returns to scale. In the next section, we add capital to the production technology. In that case, the restriction to affine production functions is not implied by constant returns.

²⁹In the absence of capital, all periods are identical, so the equilibrium allocation is time invariant and identical to the one constructed in Sections 2 to 4.

$$\sum_{t=0}^{\infty} \beta^t \mathcal{U}(c_t),\tag{59}$$

where $\beta \in (0, 1)$ is the household's discount factor.

Also, the production of goods requires not only labor but also capital:

$$y_t = \mathcal{F}(k_t, n_t). \tag{60}$$

The production function $\mathcal{F} : \Re^2_+ \to \Re_+$ maps capital and labor into output. This function is continuously differentiable, increasing in both arguments, concave, and homogeneous of degree one. Also, the Inada conditions for an interior solution are assumed to apply.

Goods can be used for both consumption and investment. When goods are used for consumption, they are exchanged in the same type of markets as those described in the previous sections. Goods used for consumption are perishable and must be sold in the same period in which they are produced. When goods are used for investment, they are perfect substitutes for one another. To save commercial costs, households use part of their own output to increase their capital stock. Therefore, households allocate a fraction of their output to be sold and another fraction to investment:

$$y_t = k_{t+1} - k_t \left(1 - \delta\right) + s_t m^s(\theta_t) \overline{q_t},\tag{61}$$

where $\delta \in (0,1)$ is the depreciation rate. The fraction destined for sale is consumed by the purchasing households.

Each period t the household is subject to constraints (5), (7) to (9), (60) and (61).³⁰ The problem of the household is to choose $\{q_{\varepsilon t}\}_{\varepsilon \in [0,1]}$, b_t , s_t , n_t , c_t , y_t and k_{t+1} to maximize (59) subject to this system of constraints. We use Lagrange multipliers $\beta^t \mu_t$, $\beta^t \lambda_t$, and $\beta^t \vartheta t$ for the resource constraint, the budget constraint, and the labor allocation constraint respectively. Then, the optimal behavior of a household is characterized by the same first-order conditions as those in Section 2, that is (10) to (13), with all variables indexed by the corresponding time subscript, and the obvious modification that output and the marginal product of labor now depend on the capital stock. Moreover, the

 $^{^{30}}$ Regarding condition (8), we could allow the clearing-house to pay or charge interest on the balance of credits and debits imposing only an intertemporal budget constraint on the households. However, since all households behave identically, all balances must be zero in equilibrium. Therefore, we can simplify the exposition by assuming that the balance of credits and debits for each household must be zero in each period.

following two conditions must hold:

$$\mu_t = \mu_{t+1} \left[\mathcal{F}_k(k_{t+1}, n_{t+1}) + 1 - \delta \right] \beta, \text{ and}$$
(62)

$$\lim_{t \to \infty} \beta^{-t} \mu_t k_t = 0. \tag{63}$$

Equation (62) states that the value of one unit of output today is equal to the present discounted value of the gross marginal product of capital, while equation (63) is a standard transversality condition.

A directed search equilibrium is a set $\{\theta_t, n_t, b_t, s_t, c_t, y_t, k_{t+1}, q_{\varepsilon t}, z_{\varepsilon t}, \psi_t, \mu_t, \lambda_t, \mathcal{Z}_t(q)\}_{t=0}^{\infty}$ satisfying conditions analogous to the ones in the definition in Section 5. This equilibrium is now described by a system of difference equations for the variables k_t and μ_t : (61) and (62) together with the time-indexed versions of conditions (1), (9), (15), (21), (30), (32), (34),^{31} (35), (37), and (60). This system together with the initial condition k_0 and the terminal condition (63) determines the equilibrium path. For all capital stocks, equations (30) and (32) still determine the pricing parameter γ_t and the congestion θ_t of buyers over sellers in the market for consumption goods. Therefore, these two variables are constant along an equilibrium path.

Qualitatively, the dynamics of capital accumulation are identical to those of the Neoclassical growth model. Capital converges monotonically to a steady state where k_t and μ_t are constant. For low capital stocks, both the marginal product of capital and the utility price of capital μ_t are high relative to the steady state. High levels of μ_t induce low consumption and high supply of labor into production, and as a result high saving. As capital is accumulated, households not only increase the fraction of output allocated for consumption, but also the fraction of labor allocated to the exchange of commodities.

In the steady state, equation (62) implies that the net marginal product of capital is equal to the subjective discount rate:

$$\mathcal{F}_k\left(\frac{k}{n},1\right) - \delta = \frac{1}{\beta} - 1.$$
(64)

When net investment is zero, analogous steps to the derivation of (38) yield:

$$\frac{\mathcal{F}_n\left(\frac{k}{n},1\right)}{\mathcal{F}\left(\frac{k}{n},1\right) - \delta\frac{k}{n}} \frac{1-n}{n} = \frac{\sigma\left(1+\gamma\right)}{1-\sigma}.$$
(65)

³¹Taking into account that the marginal product of labor depends on both capital and labor.

Equations (64) and (65) determine the steady state capital stock and labor allocated into production. The steady state values of the remaining variables (b, s, q_{ε} , and μ) are obtained from equations analogous to those in Section 4.

As in the version of the model without capital, a central planner who faces the same informational constraints as the market and who only has control of the direct revelation mechanism by which transactions are conducted cannot improve upon the allocation in a directed search equilibrium when the production technology is linear. With the existence of capital, constant returns to scale does not imply that the production technology is linear. However, with constant returns to scale the difference in welfare between a directed search equilibrium and the allocation that is attained when the central planner chooses the direct revelation trading mechanism are negligible for reasonable parameters (see the numerical example in the following subsection).

As in the Neoclassical growth model, output and capital converge to a steady growth path if the utility function \mathcal{U} is isoelastic and the efficiency of labor in the production of goods grows at a constant rate. Also, the utility function can be easily extended to include leisure or home services. In this case, the restrictions on \mathcal{U} for convergence to a steady growth path are the standard ones.

7.1 Numerical Calibration

In the context of the Neoclassical growth framework, our model of commerce can be estimated using standard economic data. In this subsection, we discuss how to identify the parameters of the model, and we provide a numerical calibration of the model. For this purpose, we assume logarithmic preferences: $\mathcal{U}(c_t) = \ln c_t$, and Cobb-Douglas matching and production technologies: $\mathcal{M}(b_t, s_t) = A_0 b_t^{\eta} s_t^{1-\eta}$ and $\mathcal{F}(k_t, e_t n_t) = A_1 k_t^{\alpha} (e_t n_t)^{1-\alpha}$.³² We also assume that the efficiency of labor in the production of goods (e_t) grows at a constant rate g.

As is standard in the Neoclassical growth model, the parameters α , β , δ , and g can be estimated using capital and labor income shares in the sector producing goods, the real return on capital, the durability of capital, and the average growth rate of output. In our numerical example in Table 2, we pick standard estimates for these parameters. The two remaining parameters of the model to estimate are the preference for diversity (σ) and the elasticity of matches with respect to the number of sellers (η).³³ These two parameters can be identified with empirical estimates of the

³²When preferences are logarithmic, the system of difference equations characterizing a directed search equilibrium is greatly simplified because then $s_t m^s (\theta_t) \overline{q_t} = (1 + \sigma \gamma)^{-1} \mu_t^{-1}$.

³³The values of the technological constants A_0 and A_1 are irrelevant for the calculations in Table 1, and affect only the units in which we measure output and hedonic consumption.

average commercial margin and the congestion in the retail sector. In a directed search equilibrium, equations (30), (33), and (41) imply:

Average commercial margin
$$\left(\frac{\lambda_t \overline{z_t} - \mu_t \overline{q_t}}{\mu_t \overline{q_t}}\right) = \frac{\sigma (1-\eta)}{1-\sigma}$$
, and (66)

Congestion
$$\left(\frac{b_t}{s_t}\right) = \frac{1+\eta}{2(1-\eta)}.$$
 (67)

Therefore, with estimates of the average commercial margin and market congestion we can solve for σ and η .

The Bureau of the Census of the United States reports that the average commercial margin for retail trade has quite stable around 0.28 during the last decade.³⁴ The Bureau of Labor Statistics measures the number of production employees in retail trade and the average weekly hours that these employees work. The product of these two measures is the empirical counterpart of s_t (465 million hours/week in 1986).³⁵ To calculate b_t , we multiply the average time spent shopping by an adult (3.4 hours/week in 1986) reported in Robinson, Andreyenkov, and Patrushev (1989, p. $84)^{36}$ by the number of shoppers in the economy (United States population 16 and over). This product measures the empirical counterpart of b_t (630 million hours/week in 1986). This implies that congestion is 1.35 ($\simeq 630/465$). Applying these estimates to the system (66) and (67), we obtain $\eta = 0.46$ and $\sigma = 0.42$.

Using the estimated parameters, Table 1 compares the directed search equilibrium and the first best allocation.³⁷ In equilibrium, households spend less time producing and selling goods and more time shopping than in the first best. Also, in equilibrium buyers leave empty handed from 27 percent of the trading meetings while they always acquire a positive amount of goods in the first best. Welfare in the two allocations is quite different. Changing from the first best balanced

³⁵We use 1986 to estimate congestion because the survey on the time individuals spend shopping that we have on hand refers to that year. The data was downloaded in July 1, 2002 from http://stats.bls.gov/data/home.htm (production employees in retail trade: 15.924 million, average hours worked by production employees: 29.2).

³⁴See http://www.census.gov/svsd/www/artstbl.html. The commercial margin varies widely by the type of business of the retailers. For example, in 2000, at the lower end, we find the commercial margins for Warehouse Clubs and Superstores (0.167), Automotive Dealers (0.175), and Gasoline Stations (0.208). At the upper end, we have the margins for Specialty Food Stores (0.419), Clothing and Footwear (0.426), and Furniture (0.441). In this paper we abstract from the reasons why different goods may trade with different margins, although this is an interesting topic for future research.

 $^{^{36}}$ This is the most recent estimate of average time allocated to shopping that we could find. We do not expect that indices of congestion vary dramatically over time.

³⁷The balance growth paths are calculated using the formulas described in the previous sections. The transitional dynamics necessary to calculate the last row are obtained using standard numerical methods.

path to the equilibrium balanced path is equivalent to a 7.94 percent drop in consumption. When the transitional costs of changing the capital stock are taken into account this percentage drops slightly to 7.87.³⁸ This large welfare cost is not easy to avoid with correcting policies. For example, regulating the price schedule in a directed search equilibrium leads only to a negligible welfare improvement equivalent to less than a 10^{-5} percent increase in consumption. (This improvement would be zero if the production technology were linear instead of Cobb-Douglas). The large difference in welfare between a directed search equilibrium and the first best allocation is due to the necessity of financing retail costs with large commercial mark-ups. This inefficiency is unavoidable in the absence of some form of lump-sum taxes which would allow to subsidize retail trade.

8 Conclusion and Extensions

Search models have been used to study decentralized markets where traders meet bilaterally. These models have been useful to analyse the labor market. Also, with the work of Kiyotaki and Wright (1989 and 1993) they have become the dominant paradigm for the theoretical microfoundations of money. Our paper uses search to capture some important features of the retail sector.

Our key assumption is that matching buyers with the sellers that carry their desired products is costly. We model this cost with the search technology of Mortensen (1982) and Pissarides (1990). We also assume that buyers are aware of the price schedules of sellers and as a result they direct their search to a subset of sellers with the most desirable combination of prices and congestion. To formalize this idea we incorporate the concept of directed search equilibrium in Peters (1991). Finally, we assume that the buyers' willingness to pay for a particular product is not observable. In this way, we extend directed search to a framework with private information.

In this framework, we study the welfare properties of a directed search allocation by comparing it to the allocation chosen by a central planner. If the planner faces the same informational constraints as the market and only has control of the direct revelation mechanism by which transactions are conducted, then a directed search equilibrium coincides with the choice of the planner when the production technology is linear. However, if the planner can use lump-sum taxes to subsidize sales, the planner can improve upon the equilibrium outcome. In fact, the planner can achieve the first best allocation by introducing a linear subsidy on sales. The first best also coincides with a directed search equilibrium if sellers can charge a flat fee to all buyers that seek to trade with them.

 $^{^{38}}$ In this comparison, we assume that both the equilibrium path and the first best path start with the same capital stock (the one in the first best).

Even if large commercial margins are unavoidable in the absence of lump-sum taxation, their existence is important for policy design. For example, large commercial margins have a profound effect on the welfare cost of sales taxes, because a zero tax rate on sales is already a large departure from the first best, so the standard result that small taxes cause negligible dead-weight-losses does not apply. Moreover, because of standard tax equivalence results, the welfare costs of income taxes must also be much larger when commercial margins are taken into account.

For simplicity, we have abstracted from the existence of money by assuming that payments can be made through a central clearing-house. In the absence of this centralized system of payments, money is a useful device that facilitates trade as in Kiyotaki and Wright (1989 and 1993). Faig (2001) introduces money in a simpler version of the present model where sellers are constrained by the set-up to make offers that consist of a single quantity-payment pair (q, z). The main complication of introducing money in the present set-up is that in equilibrium when buyers are lucky to find a good for which they have a high valuation they would like to spend more money than they carry. Because this is not possible, they are liquidity constrained. Moreover, these liquidity constraints affect equilibrium price schedules. Despite this complication, the model remains analytically tractable, but the algebraic expressions are longer and harder to interpret than in the present contribution. For this reason, we plan to study the monetary version of the present model in a separate paper.

APPENDIX

Derivation of (5)

Let q_j be the quantity purchased and ε_j be the preference shock in each match j between a buyer of the household and a seller, where j = 1, 2, ... Assume the following Dixit-Stiglitz aggregator:

$$c = \left(\lim_{J' \to \infty} \frac{1}{J'} \sum_{j=1}^{J} \varepsilon_j q_j^{1-\sigma}\right)^{\frac{1}{1-\sigma}}.$$
(68)

Here, J and J' denote the number of matches and the number of household members respectively. The goods acquired have different valuations for the household. For a given quantity purchased q_j , the household gets maximum utility when the realization of the shock in a match is $\varepsilon_j = 1$, and this utility declines linearly to 0 as ε_j decreases to 0. Because there is a countable number of potential purchases and a continuum of goods, with probability one an additional purchase brings a new good to the set of goods consumed by the household.

Since the number of household members is large, the measure of matches between the buyers of the household and a potential seller is $\lim_{J'\to\infty} J/J' = bm^b(\theta)$. Hence,

$$c = \left(bm^{b}(\theta) \lim_{J \to \infty} \frac{1}{J} \sum_{j=1}^{J} \varepsilon_{j} q_{j}^{1-\sigma} \right)^{\frac{1}{1-\sigma}}.$$
(69)

Moreover, the stochastic variables ε_j are uniformly distributed on the interval [0, 1] and are independent across meetings. Let q_{ε} denote the quantity purchased by a buyer when the realization of the preference shock is ε . The Law of Large Numbers implies the formula in (5).

Proof of Proposition 1

Let the indirect utility of a type- ε buyer be defined as

$$v_{\varepsilon} \equiv \mathcal{U}^{b}\left(q_{\varepsilon}, z_{\varepsilon}; \varepsilon\right). \tag{70}$$

Using v_{ε} , the incentive compatibility constraint (18) can be restated with the help of the following standard result (see Mas-Colell, Winston and Green, 1995, Proposition 23.D.2):

Lemma A direct revelation mechanism satisfies the incentive compatibility constraint (18) if and only if q_{ε} is non-decreasing in ε and the indirect utility function satisfies

$$v_{\varepsilon} = \int_{0}^{\varepsilon} \frac{\partial}{\partial x} \mathcal{U}^{b}(q_{x}, z_{x}; x) \, dx = \int_{0}^{\varepsilon} \frac{\psi q_{x}^{1-\sigma}}{1-\sigma} dx, \text{ for all } \varepsilon \in [0, 1].$$
(71)

Using this lemma and constraints (19) and (20) together with the definitions (14), (16), and (70), the incentive efficient direct revelation mechanisms can be characterized as the solution to the following program:

$$\max_{\{q_{\varepsilon}, v_{\varepsilon}\}_{\varepsilon \in [0,1]}} \int_0^1 \left[(1-\omega)v_{\varepsilon} + \omega \left(\frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} - \mu q_{\varepsilon} - v_{\varepsilon} \right) \right] d\varepsilon$$
(72)

subject to^{39}

$$\dot{v}_{\varepsilon} = \frac{\psi q_{\varepsilon}^{1-\sigma}}{1-\sigma},\tag{73}$$

 $q_{\varepsilon} \ge 0, \tag{74}$

$$v_{\varepsilon} \ge 0,$$
 (75)

$$\pi_{\varepsilon} \equiv \frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} - v_{\varepsilon} - \mu q_{\varepsilon} \ge 0, \tag{76}$$

$$q_{\varepsilon}$$
 is non-decreasing, (77)

$$q_0 = 0, \text{ and } v_0 = 0.$$
 (78)

Here, π_{ε} denotes the seller's surplus upon meeting a buyer of type ε . The payment in the transaction satisfies:

$$\lambda z_{\varepsilon} = \frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} - v_{\varepsilon} = \pi_{\varepsilon} + \mu q_{\varepsilon}$$
(79)

This program can be solved with a standard application of the Pontryagin's Maximum Principle. We organize the solution to the program in three steps:

(a) Constraints (73) to (78) imply that there is a $\gamma \in [0,1)$ such that the solution $\{q_{\varepsilon}, v_{\varepsilon}\}_{\varepsilon \in [0,1]}$ to program (72) obeys: $q_{\varepsilon} = 0$ and $v_{\varepsilon} = 0$ for $\varepsilon \in [0,\gamma]$, and $q_{\varepsilon} > 0$ and $v_{\varepsilon} > 0$ for $\varepsilon > \gamma$.

Constraints (77) and (78) immediately imply that there is $\gamma \in [0,1]$ such that $q_{\varepsilon} = 0$ for $\varepsilon \in [0,\gamma]$, and $q_{\varepsilon} > 0$ for $\varepsilon > \gamma$. When $q_{\varepsilon} = 0$, constraints (75) and (76) imply that $v_{\varepsilon} = 0$. With these results, constraint (73) implies that $q_{\varepsilon} > 0$ if and only if $v_{\varepsilon} > 0$. Finally, if γ were 1, the optimized value of (72) would be zero, which cannot be a solution to the maximization program because there are many feasible direct revelation mechanisms that achieve a positive value. For example, $q_{\varepsilon} = v_{\varepsilon} = 0$ if $\varepsilon < 0.5$, and $q_{\varepsilon} = \left(\frac{\psi(2\varepsilon-1)}{\mu}\right)^{\frac{1}{\sigma}}$ and $v_{\varepsilon} = \frac{1}{2}\frac{\sigma}{1-\sigma}\left[\psi\left(2\varepsilon-1\right)\right]^{\frac{1}{\sigma}}\mu^{\frac{\sigma-1}{\sigma}}$ otherwise. \Box^{40}

³⁹ \dot{v}_{ε} denotes the derivative of v with respect to ε evaluated at ε .

⁴⁰This is the optimal direct mechanism for $\omega = 1$.

(b) For $\omega \in (0.5, 1]$, the solution to program (72) subject to (73) to (78), is the following:

$$q_{\varepsilon} = \begin{cases} 0 & \text{for } 0 \le \varepsilon < \gamma \\ \left(\frac{\psi}{\mu} \frac{\varepsilon - \gamma}{1 - \gamma}\right)^{\frac{1}{\sigma}} & \text{for } \gamma \le \varepsilon \le 1 \end{cases}$$
(80)

$$v_{\varepsilon} = \begin{cases} 0 & \text{for } 0 \le \varepsilon < \gamma \\ \frac{\sigma\mu(1-\gamma)}{1-\sigma} \left(\frac{\psi}{\mu}\frac{\varepsilon-\gamma}{1-\gamma}\right)^{\frac{1}{\sigma}} & \text{for } \gamma \le \varepsilon \le 1 \end{cases}$$
(81)

where

$$\gamma = \frac{2\omega - 1}{3\omega - 1}.\tag{82}$$

Let ζ_{ε} denote the co-state variable associated with the differential equation (73). The currentvalue Hamiltonian of the program is:

$$\mathcal{H}(q_{\varepsilon}, v_{\varepsilon}, \zeta_{\varepsilon}, \varepsilon) = (1 - \omega)v_{\varepsilon} + \omega \left(\frac{\varepsilon \psi q_{\varepsilon}^{1 - \sigma}}{1 - \sigma} - \mu q_{\varepsilon} - v_{\varepsilon}\right) + \zeta_{\varepsilon} \frac{\psi q_{\varepsilon}^{1 - \sigma}}{1 - \sigma}.$$
(83)

The first-order necessary condition with respect to the control variable q_{ε} is $(\mathcal{H}_{q_{\varepsilon}} = 0)$:

$$\omega \left(\psi \varepsilon q_{\varepsilon}^{-\sigma} - \mu\right) + \zeta_{\varepsilon} \psi q_{\varepsilon}^{-\sigma} = 0.$$
(84)

The co-state variable must obey $(\dot{\zeta}_{\varepsilon} = -\mathcal{H}_{v_{\varepsilon}})$:

$$\dot{\zeta}_{\varepsilon} = -\left(1 - 2\omega\right). \tag{85}$$

Finally, the transversality implies⁴¹

$$\zeta_1 = 0. \tag{86}$$

The value of the co-state variable ζ_{ε} can be solved for explicitly using conditions (85) and (86):

$$\zeta_{\varepsilon} = (2\omega - 1) \left(\varepsilon - 1\right). \tag{87}$$

Substituting (87) into (84), solving for q_{ε} , and using (82), we obtain

$$q_{\varepsilon} = \left(\frac{\psi}{\mu}\frac{\varepsilon - \gamma}{1 - \gamma}\right)^{\frac{1}{\sigma}}.$$
(88)

Integrating (73) from γ to ε and noting that $v_{\gamma} = 0$, we obtain

$$\underline{v_{\varepsilon}} = \frac{\sigma}{1 - \sigma} \left(1 - \gamma\right) \mu q_{\varepsilon}.$$
(89)

⁴¹The transversality condition is $\zeta_1 v_1 = 0$. However, $v_1 > 0$ because of step (a).

Finally, using the definition of π_{ε} in (76), we obtain

$$\pi_{\varepsilon} = \gamma \frac{\mu q_{\varepsilon}}{1 - \sigma} \left(\sigma + \frac{1 - \varepsilon}{\varepsilon - \gamma} \right) \tag{90}$$

So far, we have disregarded constraints (74) to (78). However, for $\varepsilon \geq \gamma$, the values of q_{ε} , v_{ε} , and π_{ε} in (88), (89), and (90) are well defined and satisfy these constraints. Moreover, at $\varepsilon = \gamma$, $q_{\varepsilon} = v_{\varepsilon} = 0$. Hence, (a) implies $q_{\varepsilon} = v_{\varepsilon} = \pi_{\varepsilon} = 0$ for $\varepsilon \leq \gamma$. The Hamiltonian is strictly concave with respect to q_{ε} , and when the Hamiltonian is evaluated at the optimal choice of q_{ε} it is concave with respect to v_{ε} . Therefore, (80) to (82) characterize the unique maximum of the program. \Box (c) For $\omega \in [0, 0.5]$, the solution to program (72) subject to (73) to (78), is the following:

$$q_{\varepsilon} = \left(\frac{\psi}{\mu}\varepsilon\right)^{\frac{1}{\sigma}}$$
 and (91)

$$v_{\varepsilon} = \frac{\sigma}{1 - \sigma} \mu q_{\varepsilon} \tag{92}$$

When $\omega = 0.5$, the variable v_{ε} cancels in the objective (72) so the problem becomes separable across types. Ignoring constraints (76) and (77), the first-order conditions of the problem yield (91). Integrating (73) from 0 to ε and noting $v_0 = 0$, we obtain (92). Using the definition in (76), we obtain $\pi_{\varepsilon} = 0$ for all $\varepsilon \in [0, 1]$. The values of q_{ε} and v_{ε} in (91) and (92) satify constraints (76) and (77), so they solve program (72) subject to (73) to (78). Furthermore, the solution for $\omega = 0.5$ maximizes the expected payoff of the buyer subject to a zero expected payoff for the seller. A fortiori, given constraint (76), the same solution must apply for $\omega \in [0, 0.5)$, which assigns a lower weight to the seller in the maximized welfare function (72). \Box

Finally, Proposition 1 results from combining steps (b) and (c) together with (79).

Proof of Proposition 2

A buyer facing a price schedule γ purchases a quantity that satisfies the following specialization of (10)

$$\varepsilon \psi q_{\varepsilon}^{-\sigma} = \gamma \psi q_{\varepsilon}^{-\sigma} + (1 - \gamma) \,\mu. \tag{93}$$

Solving this equation for q_{ε} , we obtain (21). Using (24), the implied payments are given by (22). Finally, γ taking values in the interval [0, 0.5] spans all values of γ attained in (23) for $\omega \in [0, 1]$. Thus, Proposition 2 follows.

Proof of Proposition 3

Given that the equilibrium variables are recursively determined by the set of equations described in the main text. Existence and uniqueness is implied if equations (33) and (38) have a unique admissible solution. The right-hand side of (33) is a continuous and decreasing function of θ , which image spans the interval (0, 2]. The left-hand side of (33) is a continuous and non-decreasing function of θ , which is bounded away from 0 and 1. Therefore, equation (33) has a unique solution.

The left-hand side of (38) is a decreasing function of n that maps [0,1] onto $[0,\infty)$, while the right-hand side is positive and independent of n. Hence, equation (38) has a unique solution for n in the interval (0,1).

Proof of Proposition 4

Equation (47) together with the definitions of θ , m^b , and η in (1), (2), and (31), implies (48) and (51). The labor allocation constraint together with (48) and (51) yields (50). Using (44) to solve the integrals in (46) and (43), we obtain

$$f'(n) = \mathcal{M}_s(b,s) \frac{\sigma^2}{1 - \sigma^2} \left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}}, \text{ and}$$
(94)

$$f(n) = \mathcal{M}(b,s) \frac{\sigma}{1+\sigma} \left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}}.$$
(95)

The definitions (1), (2), and (31) imply

$$\frac{\mathcal{M}_s(b,s)}{\mathcal{M}(b,s)} = \frac{1 - \eta(\theta)}{s}.$$
(96)

Equation (49) results from combining (50) with (94) to (96). Finally, combining (44) and (95), we obtain (52). \blacksquare

Proof of Proposition 5

The equations that determine θ and n in a directed search equilibrium, (33) and (38)-(30), and the first best allocation, (48) and (49), can be written as follows:

$$\frac{1}{a_1 + \theta} = 1 - \eta(\theta), \text{ and}$$
(97)

$$\frac{(1-n)f'(n)}{f(n)} = \frac{\sigma}{1-\sigma}a_2,$$
(98)

where $a_1 = a_2 = 1$ in the first best, and $a_1 = 0.5$, $a_2 = (1 + \theta)/(0.5 + \theta) > 1$ in equilibrium. The Implicit Function Theorem applied to (97) and (98) yields $d\theta/da_1 < 0$ and $dn/da_2 > 0$. Hence, we obtain (54) and (53).

If f has constant elasticity α , equations (37)-(38) and (49)-(50) imply

$$\frac{s}{n} = \frac{\sigma \left[1 - \eta \left(\theta\right)\right]}{\alpha \left(1 - \sigma\right)} \tag{99}$$

for both the first best and the equilibrium. Therefore, (55) follows from (99), (53), and η being a non-increasing function of θ .

The resource constraint (7) together with f being constant returns to scale implies that

$$\frac{\overline{q}}{\overline{q}^*} = \frac{\frac{s^*}{n^*} m^s(\theta^*)}{\frac{s}{n} m^s(\theta)}.$$
(100)

Therefore, (56) follows from (55), (53), and m^s being increasing.

Proof of Proposition 6

(a) When buyer types are public information the directed search equilibrium and the first best allocations coincide.

As in Section 3, all bilateral trades must be pairwise efficient in equilibrium and the marginal rates of substitution between θ and ξ of buyers and sellers must coincide. With full information, the efficient quantities q_{ε} are calculated by maximizing the trade surplus, $\mathcal{U}^b + \mathcal{U}^s$:

$$q_{\varepsilon} = \arg \max \left(\frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} - \mu q_{\varepsilon} \right)$$
(101)

The solution to (101) is:

$$q_{\varepsilon} = \left(\frac{\psi\varepsilon}{\mu}\right)^{\frac{1}{\sigma}}.$$
(102)

The expected trade surplus is:

$$\int_0^1 \left(\frac{\varepsilon \psi q_{\varepsilon}^{1-\sigma}}{1-\sigma} - \mu q_{\varepsilon} \right) d\varepsilon = \frac{\sigma^2}{1-\sigma^2} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}}.$$
 (103)

When the market tightness is θ and buyers receive a fraction ξ of the trade surplus, the expected utilities of buyers and sellers are:

$$\mathcal{V}^{b}(\theta,\xi) = \xi m^{b}(\theta) \frac{\sigma^{2}}{1-\sigma^{2}} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}}, \text{ and}$$
(104)

$$\mathcal{V}^{s}(\theta,\xi) = (1-\xi) m^{s}(\theta) \frac{\sigma^{2}}{1-\sigma^{2}} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}}.$$
(105)

In a directed search equilibrium, the marginal rates of substitution between θ and ξ of buyers and sellers must coincide. Differentiating (104) and (105), this implies:

$$\xi = \eta(\theta). \tag{106}$$

Households allocate b and s so $\mathcal{V}^{b}(\theta,\xi) = \mathcal{V}^{s}(\theta,\xi)$. Using (104) and (105), this equality yields (48). Also, households allocate s and n so $\mathcal{V}^{s}(\theta,\alpha) = \mu f'(n)$. Using (105) and (106) this equality implies:

$$\mu f'(n) = m^s(\theta) \left[1 - \eta(\theta)\right] \frac{\sigma^2}{1 - \sigma^2} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma - 1}{\sigma}}.$$
(107)

Using (102) to calculate the average sales, the resource constraint becomes

$$f(n) = sm^{s}(\theta) \frac{\sigma}{1+\sigma} \left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}}.$$
(108)

Combining (107) and (108), and using (1), (48), and (9), we obtain (49) and (50). From (107),

$$\left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}} = \frac{f(n)}{sm^s(\theta)} \frac{1+\sigma}{\sigma},\tag{109}$$

which combined with (2) and (102) gives (52). This completes the proof of (a).

(b) If buyer types are private information but sellers charge a lump-sum fee to prospective buyers prior to the realization of ε , then the directed search equilibrium and the first best allocations coincide.

Let p be the fee a seller charges to prospective buyers. After the fee p has been paid and ε is realized, the trading game between a buyer and a seller is identical to the one in Section 3. Hence, the payment schedule net of the fee p that implements incentive efficient direct revelation mechanisms has still the functional form (24) and the quantities purchased by the buyer are given by (21). The expected utilities of the traders in a market with congestion θ , price schedule γ , and fee p are

$$\mathcal{V}^{b}(\theta,\gamma,p) = m^{b}(\theta) \left(\frac{\sigma^{2}}{1-\sigma^{2}} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}} \gamma^{2} - \lambda p \right), \text{ and}$$
(110)

$$\mathcal{V}^{s}(\theta,\gamma,p) = m^{s}(\theta) \left[\frac{2\sigma^{2}}{1-\sigma^{2}} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}} \gamma(1-\gamma) + \lambda p \right].$$
(111)

In a directed search equilibrium, the marginal rates of substitution of θ for γ , and of γ for p must be equal for buyers and sellers, so

$$\frac{\eta\left(\theta\right)-1}{\eta\left(\theta\right)} = \frac{\gamma}{1-2\gamma} \frac{\frac{2\sigma^2}{1-\sigma^2} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}} \gamma(1-\gamma) + \lambda p}{\frac{\sigma^2}{1-\sigma^2} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma-1}{\sigma}} \gamma^2 - \lambda p};$$
(112)

$$\gamma = 1. \tag{113}$$

Substituting (113) in (112) and solving for p,

$$p = \frac{1}{\lambda} \frac{\sigma^2}{1 - \sigma^2} \psi^{\frac{1}{\sigma}} \mu^{\frac{\sigma - 1}{\sigma}} \left[1 - \eta \left(\theta \right) \right].$$
(114)

Note that this fee implies $\mathcal{V}^b(\theta, \gamma, p) > 0$, so buyers are willing to pay the fee.

Households allocate b and s so $\mathcal{V}^{b}(\theta, \gamma, p) = \mathcal{V}^{s}(\theta, \gamma, p)$. Using (113) and (114), this equality yields (48). Also, households allocate s and n so $\mathcal{V}^{s}(\theta, \alpha) = \mu f'(n)$. Using (113) and (114), this

equality yields (107). To show that (49), (50) and (52) hold, we use the exact same steps used in the proof of (a). \blacksquare

Proof of Proposition 7

A type- ε buyer chooses

$$q_{\varepsilon} = \arg \max \left\{ \frac{\varepsilon \psi q^{1-\sigma}}{1-\sigma} - \lambda \mathcal{Z}(q) \right\} = \left(\frac{\psi \varepsilon}{\mu \eta} \right)^{\frac{1}{\sigma}}.$$
 (115)

Hence, the return of a seller when matched with a type- ε buyer is

$$\lambda T \mathcal{Z}(q_{\varepsilon}) - \mu q_{\varepsilon} = (T\eta - 1)\mu q_{\varepsilon} = \frac{\sigma(1 - \eta)}{1 - \sigma} \left(\frac{\psi \varepsilon}{\mu \eta}\right)^{\frac{1}{\sigma}} \mu, \tag{116}$$

where $\eta = \eta(\theta)$ and θ is the first-best level of congestion given by (48). By (48), $\eta \leq 1$. Thus, the seller's individual rationality constraint is satisfied.

When the market tightness is θ , the expected utilities of the traders are:

$$\mathcal{V}^{b}(\theta) = m^{b}(\theta)\eta \frac{\sigma^{2}}{1-\sigma^{2}} \left(\frac{\psi}{\mu\eta}\right)^{\frac{1}{\sigma}}\mu; \qquad (117)$$

$$\mathcal{V}^{s}(\theta) = m^{s}(\theta) \left(1 - \eta\right) \frac{\sigma^{2}}{1 - \sigma^{2}} \left(\frac{\psi}{\mu\eta}\right)^{\frac{1}{\sigma}} \mu.$$
(118)

The household chooses its labor allocation so that $\mathcal{V}^b(\theta) = \mathcal{V}^s(\theta) = \mu f'(n)$. The first equality implies that θ satisfies (48) so $\eta = \eta(\theta)$. The second one implies

$$f'(n) = m^{s}(\theta) \left(1 - \eta\right) \frac{\sigma^{2}}{1 - \sigma^{2}} \left(\frac{\psi}{\mu\eta}\right)^{\frac{1}{\sigma}}.$$
(119)

Using (115) to calculate the average sales, the resource constraint becomes

$$f(n) = sm^{s}(\theta) \frac{\sigma}{1+\sigma} \left(\frac{\psi}{\mu\eta}\right)^{\frac{1}{\sigma}}.$$
(120)

Combining (119) and (120), and using (1), (48), and (9), we obtain (49) and (50). From (120),

$$\left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}} = \frac{f(n)}{sm^s(\theta)} \frac{1+\sigma}{\sigma} \eta^{\frac{1}{\sigma}},\tag{121}$$

which combined with (2) and (115) gives (52).

Proof of Proposition 8

The direct revelation mechanism chosen by the planner must be incentive efficient, so we can restrict our search to price schedules in (24) for some unknown parameter γ . Substituting (21) into (5), solving for the integral, and using (2) we obtain:

$$c = \left\{ \frac{sm^s(\theta)\sigma}{1+\sigma} \left(1-\gamma\right) \left(1+\sigma\gamma\right) \right\}^{\frac{1}{1-\sigma}} \left(\frac{\psi}{\mu}\right)^{\frac{1}{\sigma}}.$$
(122)

Denoting f(n) = a + wn, and using (6), (7), (32), (35), (37), and (38), equation (122) simplifies to:

$$c = \left(\frac{2w\sigma^2}{1+\sigma}\right)^{\frac{\sigma}{1-\sigma}} (1-\sigma) \left(a+w\right)^{\frac{1}{1-\sigma}} \left[\gamma \left(1-\gamma\right) m^s(\theta)\right]^{\frac{\sigma}{1-\sigma}}.$$
(123)

Since (32) holds, the maximization of c implies:

$$(2\gamma - 1)m^{s}(\theta) + \gamma(1 - \gamma)(m^{s})'(\theta)\frac{1}{2\gamma^{2}} = 0.$$
(124)

Given (31) and (32), (124) implies (30). \blacksquare

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Table 1

NUMERICAL EXAMPLE

Production function: $\mathcal{F}(k, en) = A_1 k^{\alpha} (en)^{1-\alpha}$, $\alpha = 0.36$ Depreciation rate: 0.1 Rate of interest: r = 0.04Rate of growth of e: g = 0.018Discount factor: $\beta = (1+g)/(1+r)$ One period utility: $\ln(c)$ Preference for diversity: $\sigma = 0.42$ Matching technology: $M(b, s) = A_0 b^{\eta} s^{1-\eta}$, $\eta = 0.46$.

	Directed	First
	Search	Best
	Equilibrium	Allocation
Balanced Growth Path		
Producers (n)	0.500	0.559
Sellers (s)	0.213	0.237
Buyers (b)	0.287	0.203
Capital/Labor Ratio (k/n)	4.374	4.374
Pricing Weight Parameter (γ)	0.270	0
Quantity Purchased When $\varepsilon \geq \gamma$ (q_{ε})	$11.23(1.37\varepsilon - 0.37)^{2.38}$	$10.13\varepsilon^{2.38}$
Welfare Relative to the First Best		
(Equivalent Percentage Change in Consumption)		
Comparison Across Balanced Paths	-0.0794	0
Comparison With Same Initial Capital	-0.0787	0