

ECO 426 (Market Design) - Lecture 1

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Market: *A medium that allows buyers and sellers of a specific good or service to interact in order to facilitate an exchange.*

- markets are institutions that determine how resources are allocated
- competitive markets
 - allocation problem is solved by a price system
 - products are exchanged for a price
 - prices adjust so that supply=demand

Competitive Markets are “good” institutions for the exchange of “commodities”

- 1 the price is the only relevant variable in the economic decision to exchange
 - the identity of the counterparty is irrelevant
- 2 commodity markets are “liquid”

Sometimes prices are not “all that matter”

- job finding (allocating jobs to workers)
 - being willing to work for an employer offering a given wage does not guarantee employment
 - being willing to hire a worker demanding a given wage does not guarantee hiring
- college admissions (allocating students to colleges)
 - being willing to pay ongoing tuition does not guarantee admission
 - admitting a student does not guarantee enrollment

Sometimes prices do not “matter at all”

- marriage “market”
 - social norms (sometimes) prevent contracting on a price for the exchange
- kidney transplants (allocating kidneys from donors to patients)
 - legal (and moral) constraints prevents exchanges for valuable consideration

market liquidity concerns

Sometimes prices are the relevant variable in the transaction but markets are “illiquid”

- allocation of radio spectrum (both supply and demand side are “thin”) - also seller’s objective might be different from max profit
- trading fine art (thin supply side of van Gogh’s *Starry Night*)



Need institutions different from competitive markets to address their shortcomings

- matching markets: do not use (only) prices as allocation mechanism
- auction markets: allocation (and price formation) mechanism for “thin markets”

- Market participants are divided into two separate groups (two-sided market)
 - A set of “men” M , with a typical man $m \in M$
 - A set of “women” W , with a typical woman $w \in W$
- Allocation: each man can be matched to one woman (or stay single), and vice-versa (one-to-one matching)

- A **matching** is a collection of pairs such that:
 - each individual has one partner - (m, w) - or
 - has no partner - (m, m) - (m is “matched with self”)
- A matching can be described by a function

$$\mu : M \cup W \rightarrow M \cup W$$

such that:

- if $\mu(m) \neq m$ then $\mu(m) \in W$ (each man is either single or matched to a woman)
- if $\mu(w) \neq w$ then $\mu(w) \in M$ (each woman is either single or matched to a man)
- $\mu(\mu(x)) = x$ (if $\mu(x)$ is x 's partner, then x is $\mu(x)$'s partner)

- Each agent has a (strict) **preferences** over “acceptable” partners
 - matching with an acceptable partner is preferred to staying unmatched
- **Example:** $M = \{m_1, m_2, m_3, m_4\}$ and $W = \{w_1, w_2, w_3\}$
 - m_1 preferences $w_2 \succ w_3 \succ w_1 \succ m_1$ (or simply w_2, w_3, w_1)
 - w_1 preferences $m_1 \succ m_3 \succ w_1$ (or simply m_1, m_3)

- A matching μ is **stable** if
 - no agent is matched to an unacceptable mate - **individual rationality**
 - there is no pair of agents who would prefer to match with each other rather than their assigned partner - **no blocking**
- **Example:** $M = \{m_1, m_2\}$ and $W = \{w_1, w_2\}$
 - Both m_1 and m_2 prefer w_1 to w_2
 - Both w_1 and w_2 prefer m_1 to m_2
 - The matching $(m_1, w_2), (m_2, w_1)$ is not stable m_1 and w_1 prefer each other to their assigned partner
 - The matching $(m_1, w_1), (m_2, w_2)$ is the unique stable matching

Stable matchings examples

- **Example:** $M = \{m_1, m_2\}$ and $W = \{w_1, w_2\}$
 - m_1 prefers w_1 to w_2 and m_2 prefers w_2 to w_1
 - w_1 prefers m_1 to m_2 and w_2 prefers m_2 to m_1
 - The matching $(m_1, w_1), (m_2, w_2)$ is the unique stable matching

- **Example:** $M = \{m_1, m_2\}$ and $W = \{w_1, w_2\}$
 - m_1 prefers w_1 to w_2 and m_2 prefers w_2 to w_1
 - w_1 prefers m_2 to m_1 and w_2 prefers m_1 to m_2
 - The matching $(m_1, w_1), (m_2, w_2)$ is stable
 - The matching $(m_1, w_2), (m_2, w_1)$ is also stable
 - Note both men prefer the first matching and the opposite is true for the women

- **Question:** can we always find a stable matching?

Gale and Shapley '62 - deferred acceptance algorithm

- Round 1
 - Each man proposes to his most preferred woman
 - Each woman reject all but the most preferred proposal received
 - Most preferred proposal is "held" not accepted
- Round 2
 - Each man rejected in the previous round proposes to the most preferred woman who has not yet rejected him
 - Each woman rejects all but the most preferred among new and held proposals
 - Most preferred proposal is "held" not accepted
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- Round T - no proposal is rejected
 - algorithm ends
 - each woman is matched to the currently held proposal

Gale and Shapley '62 - deferred acceptance algorithm

Example $M = \{m_1, m_2, m_3\}$ and $W = \{w_1, w_2, w_3\}$

m_1	w_1	w_2	w_3		w_1	m_1	m_2	m_3
m_2	w_3	w_1			w_2	m_1	m_2	m_3
m_3	w_1	w_3	w_2		w_3	m_3	m_1	m_2

m_1	m_2	m_3
w_1	w_3	w_1
w_1	w_3	\emptyset
w_3		
w_1	\emptyset	w_3
w_1		
w_1	\emptyset	w_3
m_2		
w_1	m_2	w_3
w_1	m_2	w_3

Existence of a stable matching

The deferred acceptance (DA) algorithm ends in finitely many rounds

- No man ever proposes twice to the same woman
- The outcome of the DA algorithm is a matching

Theorem. The outcome of the DA algorithm is a stable matching.

- The proposal “held” by each woman improves (weakly) in each round
 - Each woman is matched to the most desirable man who has ever proposed to her
- Every time he makes a new proposal, a man proposes to the next most desirable woman
 - Each man has proposed to, and has been rejected by, all women more desirable than his match
- For each man m , every woman more desirable than his match prefers her current match to m .

- Existence of a stable matching relies on the two-sided nature of the market
- Example (the roommate problem) Three students A, B, C can form a pair to share a two-people room
 - A prefers sharing with B to sharing with C to not sharing
 - B prefers sharing with C to sharing with A to not sharing
 - C prefers sharing with A to sharing with B to not sharing
 - there is no stable matching

- When multiple stable matching exist, the DA algorithm with man proposing yields a different outcome from the DA algorithm with women proposing

multiple stable matching - conflicting preferences

Example $M = \{m_1, m_2, m_3\}$ and $W = \{w_1, w_2, w_3\}$

m_1	w_1	w_2	w_3
m_2	w_1	w_2	w_3
m_3	w_1	w_3	w_2

w_1	m_1	m_2	m_3
w_2	m_1	m_3	m_2
w_3	m_1	m_2	m_3

DA m

m_1	m_2	m_3
w_1	w_1	w_1
w_1	\emptyset	\emptyset
w_1	w_2	w_3
w_1	w_2	w_3
w_1	w_2	w_3

DA w

w_1	w_2	w_3
m_1	m_1	m_1
m_1	\emptyset	\emptyset
m_1	m_3	m_2
m_1	m_3	m_2
m_1	m_3	m_2

multiple stable matching - conflicting preferences

- **Example** $M = \{m_1, m_2, m_3\}$ and $W = \{w_1, w_2, w_3\}$

m_1	w_1	w_2	w_3	w_1	m_1	m_2	m_3
m_2	w_1	w_2	w_3	w_2	m_1	m_3	m_2
m_3	w_1	w_3	w_2	w_3	m_1	m_2	m_3

- DA men proposing: $(m_1, w_1), (m_2, w_2), (m_3, w_3)$
- DA women proposing $(m_1, w_1), (m_2, w_3), (m_3, w_2)$
- all men prefer (weakly) the DA men proposing outcome
- all women prefer (weakly) the DA women proposing outcome
- **Theorem** For two stable matchings μ, μ' , all men (weakly) prefer μ if and only if all women (weakly) prefer μ' .
- The DA algorithm with men proposing yields the best stable matching for men and the worst stable matching for women.